Computer Science Jerzy Świątek Systems Modelling and Analysis

Choose yourself and new technologies

#### L.17. Modeling of complex of operation systems









#### Identification of complexes of operations with restricted measurements possibilities

- Description of complexes of operations.
- Identification of complexes of operations.
  - unlimited measurement possibilities,
  - limited possibilities of measurement of operations execution time,
  - limited possibilities of measurement of operations execution time and size of task or amount of resources allocated to an operation.
- Final remarks
- References

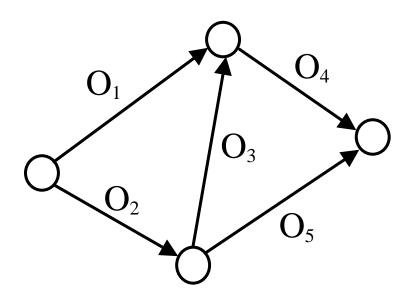








## Description of complex of operations











# Description of complex of operations

 $O_1, O_2, \ldots, O_M$  – elementary static operations

For m-th operation description is given:

$$T_m = F_m(u_m, a_m), \quad m = 1, 2, \dots, M,$$

where  $T_m \ge 0$  *m*-th operation completion time,

 $u_m - s_m$ -dimensional vector of m-th operation's inputs:  $u_m \in U_m \subseteq \mathcal{R}^{+s_m}$ ,  $a_m - r_m$ -dimensional vector of parameters:  $a_m \in A_m \subseteq \mathcal{R}^{r_m}$ ,  $F_m$  – known function:  $F_m : U_m \times A_m \to \mathcal{R}^+$ .









# Description of complex of operations

Coordinates of the vector  $u_m$  stand for amount of resources or size of task for *m*-th operation.

#### **Resources:**

 $F_m$  – is nonincreasing function with respect to all of the vector  $u_m$  coordinates For each  $a_m$  we have:

$$F_m(0_m,a_m)=\infty.$$

#### Tasks:

 $F_m$  – is nondecreasing function with respect to all of the vector  $u_m$  coordinates For each  $a_m$  we have:

$$F_m(0_m,a_m)=0.$$









# Description of complex of operations

**Structure of the system** is described by the following graph:

$$G \subset \{1, 2, \ldots, M\} \times \{1, 2, \ldots, M\}$$

If  $(m, n) \in G$  then the *m*-th operation is performed just after the *n*-th operation ends up.

The whole system **completion time**:  $T = H(T_1, T_2, ..., T_M)$ ,

where H – function determining the complex of operations completion time, dependent on the complex of operations structure.

$$T = H(F_1(u_1, a_1), F_2(u_2, a_2), \dots, F_M(u_M, a_M)) = F(u_1, u_2, \dots, u_M, a_1, a_2, \dots, a_M)$$

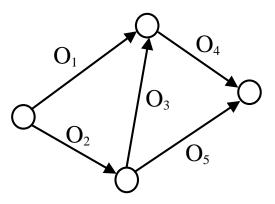








# Identification of complex of operations



 $T_m = F_m(u_m, a_m), \quad m = 1, 2, \dots, M, \quad T = H(T_1, T_2, \dots, T_M)$ 

H – function determining the total runtime of complex of operation

 $F_1, F_2, \dots, F_M$  – known functions  $a_1, a_2, \dots, a_M$  – unknown parameters









## Unlimited measurement possibilities

Available measurements:  $T_m(n), u_m(n), m = 1, 2, ..., M$ ,

where:  $T_m(n)$  – measurement of *m*-th operation completion time for resource  $u_m(n)$ 

For each operation:

$$T_m(n) = F_m(u_m(n), a_m), \quad n = 1, 2, ..., N.$$

Solving this system of equations with respect  $a_m$  results in identification algorithm for m-th operation,



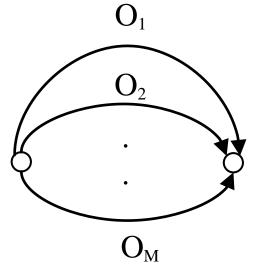






## Unlimited measurement possibilities

• Example – tasks allocation for complex of parallel operations



Complex of parallel operations









## Unlimited measurement possibilities

• Example – tasks allocation for complex of parallel operations

For the complex of operations we have description:

$$T_m = a_m u_m, \quad a_m > 0, \ u_m \ge 0, \ m = 1, 2, \dots, M$$

The total size of all taks is u.

Solution of tasks allocation problem should satisfy the following constraints:

$$\mathcal{D}_{u} = \left\{ u_{m} \geq 0, \ m = 1, 2, \dots, M; \ \sum_{m=1}^{M} u_{m} = u \right\}$$









## Unlimited measurement possibilities

• Example – tasks allocation for complex of parallel operations

For each operation we have:  $T_m(n) = a_m u_m(n)$ 

Note, that for m-th operation one measurement is enough.

Parameter of m-th operation's description we evaluate as:

$$a_m = \frac{T_m(n)}{u_m(n)}, \quad m = 1, 2, \dots, M$$





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Available measurements: $T(n), u_1(n), u_2(n), \dots, u_M(n), \quad n = 1, 2, \dots, N,$ where:T(n) – measurement of m-th operation completion time for resource  $u_m(n)$ N – number of experiment's repetitions $m = 1, 2, \dots, M$ 

Observed completion time of complex of operations for measured data is given by:

$$T(n) = F(u_1(n), u_2(n), \dots, u_M(n), a_1, a_2, \dots, a_M), \quad n = 1, 2, \dots, N$$

Solving this system of equations with respect  $a_1, a_2, ..., a_M$  results in identification algorithm.

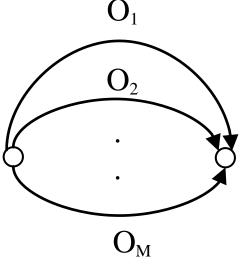








• Example – tasks allocation for complex of parallel operations



Complex of parallel operations









• Example – tasks allocation for complex of parallel operations

**Total completion time** for the whole complex of operations is given by:

$$T = \max_{1 \le m \le M} \{a_m u_m\}$$

For measurement data we have:

$$T(n) = \max_{1 \le m \le M} \{a_m u_m(n)\}, \quad n = 1, 2, \dots, N.$$

Solving this system of equations with respect  $a_1, a_2, ..., a_M$  results in identification algorithm.









• Example – tasks allocation for complex of parallel operations

**Total completion time** for the whole complex of operations is given by:

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For measurement data we have:

$$T(n) = \max_{1 \le m \le M} \{a_m u_m(n)\}, \quad n = 1, 2, \dots, N.$$

Solving this system of equations with respect  $a_1, a_2, \ldots, a_M$  results in identification algorithm.









• Example – tasks allocation for complex of parallel operations

For n -th run of complex of operations we allocate all resources or tasks to a single operation:

$$u_m(m) = u(m), \quad u_m(n) = 0, \quad n = 1, 2, \dots, M, \ n \neq m$$

For such an experiment we have:

$$T(m) = a_m u_m(m), \quad m = 1, 2, ..., M.$$

Solving this system of equations with respect  $a_1, a_2, ..., a_M$  results in identification algorithm in the form:

$$a_m = \frac{T(m)}{u_m(m)}, \quad m = 1, 2, \dots, M.$$









• Example – tasks allocation for complex of parallel operations

We allocate all resources or tasks uniformly to each operation:

$$u_1(n) = u_2(n) = \dots = u_M(n) = \overline{u}(n) = \frac{u(n)}{M}$$

For such an experiment we have:

$$T(n) = \overline{u}(n) \max_{1 \le m \le M} \{a_m\}, \quad n = 1, 2, ..., N.$$

Solution of the system of equation is not unique with respect  $a_1, a_2, ..., a_M$ . We are only able to work out a parameter, which is a function of parameters  $a_1, a_2, ..., a_M$ : T(n)

$$\max_{1 \le m \le M} \{a_m\} = \frac{T(n)}{\overline{u}(n)}.$$









Available measurements:  $T^*(n), u(n), n = 1, 2, ..., N$ ,

where:  $T^*(n)$  – optimal completion time for the maximal size of task or the global resource  $u_m(n)$ 

Following assumption about optimal task or resources allocation, we take the following allocation into consideration:

$$u_1^*, u_2^*, \dots, u_M^*.$$

For such allocation the completion time is minimal:

$$F(u_1^*, u_2^*, \dots, u_M^*, a_1, a_2, \dots, a_M) = \min_{(u_1, u_2, \dots, u_M) \in \mathcal{D}_u} F(u_1, u_2, \dots, u_M, a_1, a_2, \dots, a_M).$$









Solution of the problem results in optimal algorithms of allocation:

$$u_m^* = G_m(u, a_1, a_2, \dots, a_M), \quad m = 1, 2, \dots, M.$$

Optimal completion time for complex of operations:

$$T^* = F(u_1^*, u_2^*, \dots, u_M^*, a_1, a_2, \dots, a_M) = F(G_1(u, a_1, a_2, \dots, a_M), \dots, G_M(u, a_1, a_2, \dots, a_M), a_1, a_2, \dots, a_M) = \widetilde{F}(u, a_1, a_2, \dots, a_M)$$

For observed measurement data we may propose the following system of equations:

$$T^*(n) = \widetilde{F}(u(n), a_1, a_2, \dots, a_M), \quad n = 1, 2, \dots, N.$$

Solving this system of equations with respect  $a_1, a_2, \ldots, a_M$  results in identification algorithm.

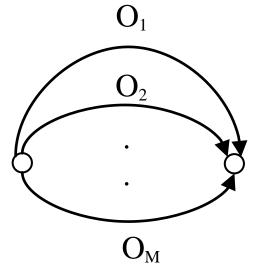








• Example – tasks allocation for complex of parallel operations











• Example – tasks allocation for complex of parallel operations

The completion time of complex of operations is optimal as long as all operations are completed at the same moment:

$$T^* = T_1^* = T_2^* = \ldots = T_M^*.$$

Taking description of operations and constraints, optimal task allocation satisfies the following system of equations:

$$T^* = a_m u_m, \quad m = 1, 2, \dots, M, \quad \sum_{m=1}^M u_m = u.$$







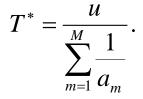


Example – tasks allocation for complex of parallel operations

Solving the system of equations with respect  $u_1, u_2, ..., u_M$  results in **optimal allocation algorithm** in the form:

$$u_m^* = \frac{u}{a_m \sum_{m=1}^M \frac{1}{a_m}}, \quad m = 1, 2, \dots, M.$$

**Optimal completion time** is expressed by the formula:







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• Example – tasks allocation for complex of parallel operations

For measurement data we have: 
$$T^*(n) = \frac{u(n)}{\sum_{m=1}^{M} \frac{1}{a_m}}, \quad n = 1, 2, ..., N.$$

Solution of the system of equation is not unique with respect  $a_1, a_2, ..., a_M$ . We are only able to work out a parameter, which is a function of

parameters  $a_1, a_2, \dots, a_M$ :  $\frac{u(n)}{\sum_{m=1}^M \frac{1}{a_m}} = \frac{T^*}{u(n)}, \quad n = 1, 2, \dots, N.$ 

Obtained value may be treated as parameter of operation equivalent to the whole complex of operations.









#### Final remarks

- Identification of complexes of operations.
  - unlimited measurement possibilities,
  - limited possibilities of measurement of operations execution time,
  - limited possibilities of measurement of operations execution time and size of task or amount of resources allocated to an operation.
- The problem of separability.









#### References

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#### Thank you for attention

